Computational optimal transport: recent speed-ups and applications

Jean Feydy HeKA team, Inria Paris Inserm, Université Paris-Cité

11th of July, 2024 SciML 2024, Université de Strasbourg

Who am I?

Background in **mathematics** and **data sciences**:

- **2012–2016** ENS Paris, mathematics.
- **2014–2015** M2 mathematics, vision, learning at ENS Cachan.
- **2016–2019** PhD thesis in **medical imaging** with Alain Trouvé at ENS Cachan.
- **2019–2021 Geometric deep learning** with Michael Bronstein at Imperial College.
 - **2021+ Medical data analysis** in the HeKA INRIA team (Paris).

HeKA: a translational research team for public health

Hôpitaux

Inria Inserm

Universités



My main motivation

Develop **robust and efficient** software that **stimulates other researchers**:

- 1. Speed up **geometric machine learning** on GPUs:
 - ⇒ **pyKeOps** library for distance and kernel matrices, 600k+ downloads.
- 2. Scale up **pharmacovigilance** to the full French population:
 - ⇒ **survivalGPU**, a fast re-implementation of the R survival package.
- 3. Ease access to modern statistical **shape analysis**:
 - ⇒ GeomLoss, truly scalable optimal transport in Python.
 - ⇒ **scikit-shapes**, alpha release now available.

Today's talk - assuming that you would enjoy some nice simulations

- 1. A quick heads up on **fast geometric methods**.
- 2. Efficient discrete optimal transport solvers.
- 3. New applications for systems of **incompressible particles**.

How to code a N-body simulation?

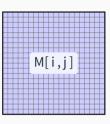
Scientific computing libraries represent most objects as tensors

Context. Constrained **memory accesses** on the GPU:

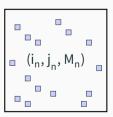
- Long access times to the registers penalize the use of large dense arrays.
- Hard-wired contiguous memory accesses penalize the use of sparse matrices.

Challenge. In order to reach optimal run times:

- **Restrict** ourselves to operations that are supported by the constructor: convolutions, FFT, etc.
- Develop new routines from scratch in C++/CUDA (FAISS, KPConv...): several months of work.



Dense array



Sparse matrix

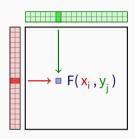
The KeOps library: efficient support for symbolic matrices

Solution. KeOps-www.kernel-operations.io:

- For PyTorch, NumPy, Matlab and R, on **CPU and GPU**.
- Automatic differentiation.
- Just-in-time compilation of optimized C++ schemes, triggered for every new reduction: sum, min, etc.

If the formula "F" is simple (\leqslant 100 arithmetic operations): "100k \times 100k" computation \rightarrow 10ms – 100ms, "1M \times 1M" computation \rightarrow 1s – 10s.

Hardware ceiling of 10^{12} operations/s. \times **10 to** \times **100 speed-up** vs standard GPU implementations for a wide range of problems.



Symbolic matrix Formula + data

- Distances d(x_i,y_j).
- Kernel k(x_i,y_i).
- Numerous transforms.

A first example: efficient nearest neighbor search in dimension 50

Create large point clouds using **standard PyTorch syntax**:

import torch N, M, D = 10**6, 10**6, 50 x = torch.rand(N, 1, D).cuda() # (1M, 1, 50) array y = torch.rand(1, M, D).cuda() # (1, 1M, 50) array

Turn dense arrays into symbolic matrices:

```
from pykeops.torch import LazyTensor
x_i, y_j = LazyTensor(x), LazyTensor(y)
```

Create a large **symbolic matrix** of squared distances:

```
D_{ij} = ((x_i - y_j) ** 2).sum(dim=2) # (1M, 1M) symbolic
```

Use an .argmin() **reduction** to perform a nearest neighbor query:

```
indices_i = D_ij.argmin(dim=1) # -> standard torch tensor
```

The KeOps library combines performance with flexibility

Script of the previous slide = efficient nearest neighbor query, on par with the bruteforce CUDA scheme of the FAISS library... And can be used with any metric!

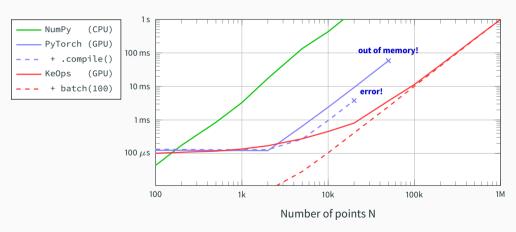
```
D_ij = ((x_i - x_j) ** 2).sum(dim=2)  # Euclidean
M_ij = (x_i - x_j).abs().sum(dim=2)  # Manhattan
C_ij = 1 - (x_i | x_j)  # Cosine
H_ij = D_ij / (x_i[...,0] * x_j[...,0])  # Hyperbolic
```

KeOps supports arbitrary **formulas** and **variables** with:

- Reductions: sum, log-sum-exp, K-min, matrix-vector product, etc.
- **Operations:** +, \times , sqrt, exp, neural networks, etc.
- Advanced schemes: batch processing, block sparsity, etc.
- Automatic differentiation: seamless integration with PyTorch.

KeOps lets users work with millions of points at a time

Benchmark of a Gaussian **convolution** $a_i \leftarrow \sum_{j=1}^N \exp(-\|x_i - y_j\|_{\mathbb{R}^3}^2) \, b_j$ between **clouds of N 3D points** on a A100 GPU.



Yet another ML compiler?

Many impressive tools out there (Taichi, Numba, Triton, Halide...):

- Focus on **generality** (software + hardware).
- Increasingly easy to use via e.g. PyTorch 2.0.

KeOps fills a different niche (a bit like cuFFT, FFTW...):

- Focus on a **single major bottleneck**: geometric interactions.
- Agnostic with respect to Euclidean / non-Euclidean formulas.
- Fully compatible with PyTorch, NumPy, R.
- Can actually be used by mathematicians.

KeOps is a **bridge** between geometers (with a maths background) and compiler experts (with a CS background).

Optimal transport?

Optimal transport (OT) generalizes sorting to spaces of dimension ${\sf D}>1$

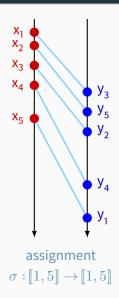
If $A=(x_1,\ldots,x_N)$ and $B=(y_1,\ldots,y_N)$ are two clouds of N points in \mathbb{R}^D , we define:

$$\mathsf{OT}(\mathsf{A},\mathsf{B}) \ = \ \min_{\sigma \in \mathcal{S}_\mathsf{N}} \ \frac{1}{\mathsf{2N}} \sum_{\mathsf{i}=\mathsf{1}}^\mathsf{N} \| \, \mathbf{x}_{\mathsf{i}} - \mathbf{y}_{\sigma(\mathsf{i})} \|^2$$

Generalizes **sorting** to metric spaces.

Linear problem on the permutation matrix P:

$$\begin{split} \mathsf{OT}(\mathsf{A},\mathsf{B}) \; &= \; \min_{\mathsf{P} \in \mathbb{R}^{\mathsf{N} \times \mathsf{N}}} \; \frac{1}{2\mathsf{N}} \sum_{i,\,j=1}^{\mathsf{N}} \mathsf{P}_{i,j} \cdot \| \, \mathbf{x}_i - \mathbf{y}_j \|^2 \; , \\ \text{s.t.} \quad \mathsf{P}_{i,j} \; &\geqslant \; 0 \qquad \underbrace{\sum_{j} \mathsf{P}_{i,j} \; = \; 1}_{\text{Each source point...}} \quad \underbrace{\sum_{i} \mathsf{P}_{i,j} \; = \; 1}_{\text{is transported onto the target.}} \end{split}$$



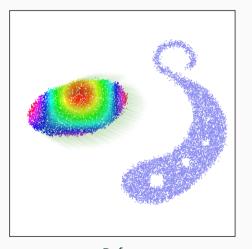
Practical use

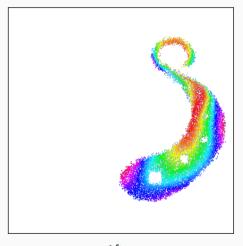
Alternatively, we understand OT as:

- Nearest neighbor **projection** + **incompressibility** constraint.
- Fundamental example of **linear optimization** over the transport plan $P_{i,j}$.

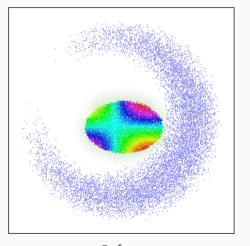
This theory induces two main quantities:

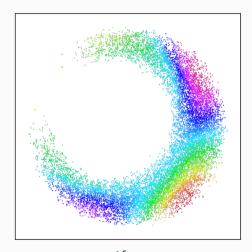
- The transport plan $\mathsf{P}_{i,j} \simeq$ the optimal mapping $\pmb{x_i} \mapsto y_{\sigma(i)}.$
- The "Wasserstein" distance $\sqrt{\mathsf{OT}(\mathsf{A},\mathsf{B})}$.



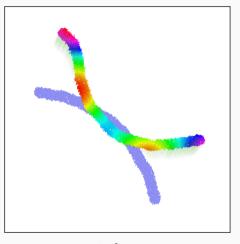


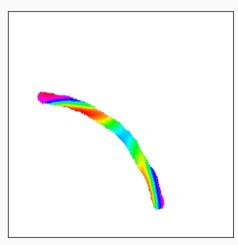
Before After ₁₄





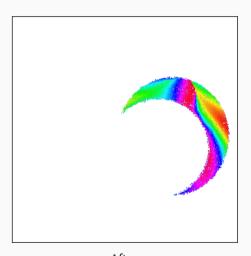
Before After ₁₄





Before After



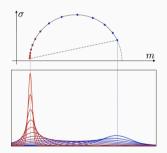


Before After 14

OT induces a geometry-aware distance between probability distributions [PC18]

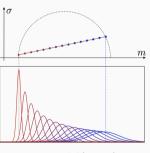
$$\text{Gauss map} \quad \mathcal{N}: (m,\sigma) \in \mathbb{R} \times \mathbb{R}_{\geq 0} \quad \mapsto \quad \mathcal{N}(m,\sigma) \in \mathbb{P}(\mathbb{R}).$$

If the space of **probability distributions** $\mathbb{P}(\mathbb{R})$ is endowed with a given metric, what is the "pull-back" geometry on the space of **parameters** (m, σ) ?



Fisher-Rao (\simeq relative entropy) on $\mathcal{N}(m,\sigma)$

ightarrow Hyperbolic **Poincaré** metric on (m,σ) .



OT on $\mathcal{N}(m,\sigma)$

 \rightarrow Flat **Euclidean** metric on (m, σ) .

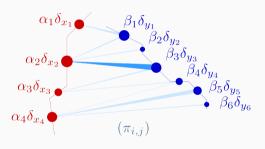
How to solve the OT problem?

Duality: central planning with NM variables \simeq outsourcing with N + M variables

$$\begin{array}{lll} \mathsf{OT}(\mathsf{A},\mathsf{B}) \ = \ \min_{\pi} \ \langle \, \pi \,,\, \mathsf{C} \, \rangle, \ \ \mathsf{with} \ \mathsf{C}(\pmb{x_i},y_j) \ = \ \frac{1}{p} \|\pmb{x_i} - y_j\|^p & \longrightarrow \ \ \mathsf{Assignment} \\ & \mathsf{s.t.} \ \pi \, \geqslant \, 0, \quad \pi \, \mathbf{1} \ = \ \mathsf{A}, \quad \pi^\mathsf{T} \mathbf{1} \ = \ \mathsf{B} \end{array}$$





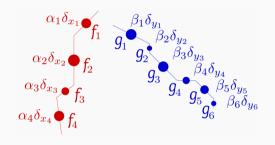


Duality: central planning with NM variables \simeq outsourcing with N + M variables

$$\begin{split} \mathsf{OT}(\mathsf{A},\mathsf{B}) \ = \ \min_{\pi} \ \langle \, \pi \,,\, \mathsf{C} \, \rangle, \ \ \mathsf{with} \ \mathsf{C}(\pmb{x_i},y_j) \ = \ \tfrac{1}{p} \|\pmb{x_i} - y_j\|^p \qquad \longrightarrow \ \ \mathsf{Assignment} \\ \mathsf{s.t.} \ \pi \ \geqslant \ 0, \quad \pi \, \mathbf{1} \ = \ \mathsf{A}, \quad \pi^\mathsf{T} \mathbf{1} \ = \ \mathsf{B} \end{split}$$



$$\sum_{i,j} \pi_{i,j} \, \mathsf{C}(\pmb{x_i}, y_j)$$



$$\max_{f,g} \qquad \langle \mathsf{A}\,,\,f \rangle \,+\, \langle \,\mathsf{B}\,,\,g \, \rangle$$
 s.t. $f(x_i)\,+\,g(y_j)\,\leqslant\, \mathsf{C}(x_i,y_j),$



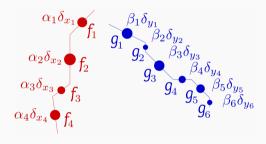
$$rac{1}{n} lpha_i f_i + \sum_j eta_j g_j \
ightarrow ext{ FedEx}$$

Duality: central planning with NM variables \simeq outsourcing with N + M variables

$$\begin{split} \mathsf{OT}(\mathsf{A},\mathsf{B}) \ = \ \min_{\pi} \ \langle \, \pi \,,\, \mathsf{C} \, \rangle, \ \ \mathsf{with} \, \mathsf{C}(\pmb{x_i},y_j) \ = \ \tfrac{1}{p} \|\pmb{x_i} - y_j\|^p \qquad \longrightarrow \ \ \mathsf{Assignment} \\ \mathsf{s.t.} \ \pi \ \geqslant \ 0, \quad \pi \, \mathbf{1} \ = \ \mathsf{A}, \quad \pi^\mathsf{T} \mathbf{1} \ = \ \mathsf{B} \end{split}$$



$$\sum_{i,j} \pi_{i,j} \, \mathsf{C}(oldsymbol{x_i}, y_j)$$



$$= \max_{f,\,g} \qquad \langle \mathsf{A}\,,\,f
angle \, + \, \langle \,\mathsf{B}\,,\,g \,
angle$$
 s.t. $f(x_i) \,+\, g(y_i) \,\leqslant\, \mathsf{C}(x_i,y_i),$



Being too greedy... doesn't work!

$$\begin{split} \text{OT}(\alpha,\beta) = & \max_{\substack{(f_i) \in \mathbb{R}^{\text{N}} \\ (g_j) \in \mathbb{R}^{\text{M}}}} \sum_{i=1}^{\text{N}} \alpha_i f_i + \sum_{j=1}^{\text{M}} \beta_j g_j \\ & \text{s.t. } \forall i,j, \ f_i + g_j \ \leqslant \ \mathbf{C}(x_i,y_j) \end{split}$$

Algorithm 3.1: Naive greedy algorithm

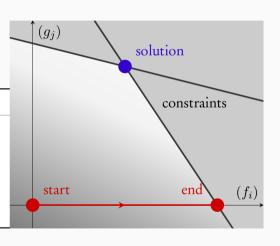
1:
$$f_i, g_j \leftarrow \mathbf{0}_{\mathbb{R}^{\mathrm{N}}}, \mathbf{0}_{\mathbb{R}^{\mathrm{M}}}$$

2: repeat

3:
$$f_i \leftarrow \min_{j=1}^{M} \left[\mathbf{C}(x_i, y_j) - g_j \right]$$

4:
$$g_j \leftarrow \min_{i=1}^{N} \left[\mathbf{C}(x_i, y_j) - f_i \right]$$

- 5: until convergence.
- 6: return f_i , g_j



The auction algorithm: take it easy with a slackness $\varepsilon>0$

$$\begin{aligned} \text{OT}(\alpha,\beta) &= & \max_{\substack{(f_i) \in \mathbb{R}^{\mathrm{N}} \\ (g_j) \in \mathbb{R}^{\mathrm{M}}}} \sum_{i=1}^{\mathrm{N}} \alpha_i f_i + \sum_{j=1}^{\mathrm{M}} \beta_j g_j \\ &\text{s.t. } \forall i,j, \ f_i + g_j \ \leqslant \ \mathbf{C}(x_i,y_j) \end{aligned}$$

Algorithm 3.2: Pseudo-auction algorithm

1:
$$f_i, g_j \leftarrow \mathbf{0}_{\mathbb{R}^{\mathrm{N}}}, \mathbf{0}_{\mathbb{R}^{\mathrm{M}}}$$

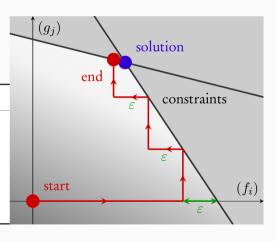
2: repeat

3:
$$f_i \leftarrow \min_{j=1}^{M} \left[\mathbf{C}(x_i, y_j) - g_j \right] - \varepsilon$$

4:
$$g_j \leftarrow \min_{i=1}^{N} [\mathbf{C}(x_i, y_j) - f_i]$$

5: **until** $\forall i, \exists j, f_i + g_j \geqslant \mathbf{C}(x_i, y_j) - \varepsilon$.

6: return f_i , g_j



The Sinkhorn algorithm: use a softmin, get a well-defined optimum

$$\begin{aligned} \text{OT}(\alpha, \beta) &= & \max_{\substack{(f_i) \in \mathbb{R}^N \\ (g_j) \in \mathbb{R}^M}} \sum_{i=1}^N \alpha_i f_i + \sum_{j=1}^M \beta_j g_j \\ &- \varepsilon \log \left\langle \alpha_i \otimes \beta_j, \exp \frac{1}{\varepsilon} [f_i \oplus g_j - \mathbf{C}_{ij}] \right\rangle \end{aligned}$$

Algorithm 3.3: Sinkhorn or "soft-auction" algorithm

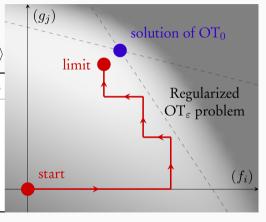
1:
$$f_i, g_j \leftarrow \mathbf{0}_{\mathbb{R}^{\mathrm{N}}}, \mathbf{0}_{\mathbb{R}^{\mathrm{M}}}$$

2: repeat

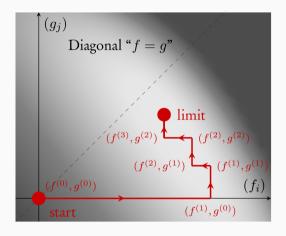
3:
$$f_i \leftarrow -\varepsilon \log \sum_{j=1}^{M} \beta_j \exp \frac{1}{\varepsilon} [g_j - \mathbf{C}(x_i, y_j)]$$

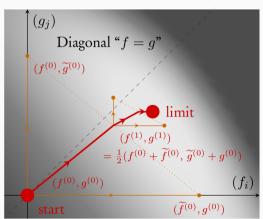
4:
$$g_j \leftarrow -\varepsilon \log \sum_{i=1}^{N} \alpha_i \exp \frac{1}{\varepsilon} [f_i - \mathbf{C}(x_i, y_j)]$$

- 5: **until** convergence up to a set tolerance.
- 6: return f_i , g_j



The symmetric Sinkhorn algorithm: stay close to the diagonal if A \simeq B





Remark 1: a streamlined algorithm

One key operation – the soft, **weighted distance transform**:

$$\forall i \in [1, \mathbf{N}], \ f(x_i) \leftarrow \min_{y \sim \beta} \left[\mathsf{C}(x_i, y) - g(y) \right] = -\varepsilon \log \sum_{j=1}^{\mathsf{M}} \beta_j \exp \tfrac{1}{\varepsilon} \big[g_j - \mathsf{C}(x_i, y_j) \big] \ .$$

Similar to the chamfer distance transform, convolution with a Gaussian kernel... Fast implementations with **pyKeOps**:

- If $C(x_i, y_j)$ is a closed formula: **bruteforce** scales to N, M \simeq 100k in 10ms on a GPU.
- If A and B have a low-dimensional support:
 use a clustering and truncation strategy to get a x10 speed-up.
- If A and B are supported on a 2D or 3D grid and $C(x_i, y_j) = \frac{1}{2} ||x_i y_j||^2$: use a **separable** distance transform to get a second x10 speed-up. (N.B.: FFTs run into numerical accuracy issues.)

Remark 2: annealing works!

The **Auction/Sinkhorn** algorithms:

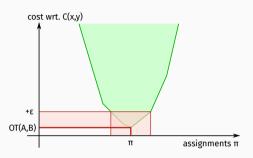
- Improve the dual cost by at least ε at each (early) step.
- Reach an ε -optimal solution with $(\max C) / \varepsilon$ steps.

Simple heuristic: run the optimization with **decreasing values** of ε .

 ε -scaling

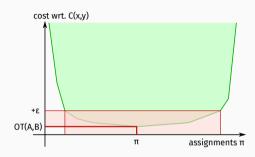
- simulated annealing
- = multiscale strategy
- = divide and conquer

Remark 3: the curse of dimensionality



In low dimension:

- ||x y|| takes large and small values.
- The OT objective is **peaky** wrt. π .
- ε -optimal solutions are **useful**.
- OT(discrete samples) ≃
 OT(underlying distributions)



In **high dimension**:

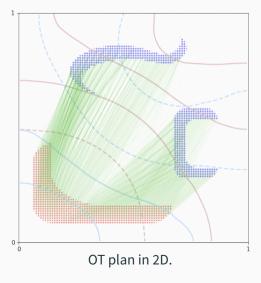
- $\|x-y\|$ gets closer to a constant.
- The OT objective is **flat** wrt. π .
- ε -optimal solutions are **random**.
- OT(discrete samples) \neq OT(underlying distributions)

To recap 80+ years of work...

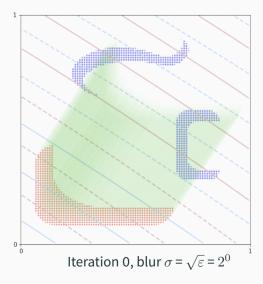
Key dates for discrete optimal transport with N points:

- [Kan42]: **Dual** problem of Kantorovitch.
- [Kuh55]: **Hungarian** methods in $O(N^3)$.
- [Ber79]: **Auction** algorithm in $O(N^2)$.
- [KY94]: **SoftAssign** = Sinkhorn + simulated annealing, in $O(N^2)$.
- [GRL+98, CR00]: **Robust Point Matching** = Sinkhorn as a loss.
- [Cut13]: Start of the GPU era.
- [Mér11, Lév15, Sch19]: **multi-scale** solvers in $O(N \log N)$.
- Solution, today: Multiscale Sinkhorn algorithm, on the GPU.
 - \Longrightarrow Generalized **QuickSort** algorithm.

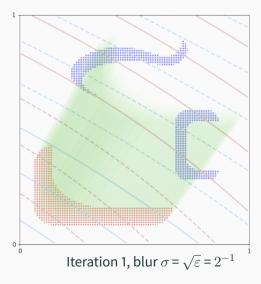
Visualizing F , G and the Brenier map $\nabla F(x_i) = -\frac{1}{\alpha_i}\partial_{x_i}\mathrm{OT}(\alpha,\beta)$



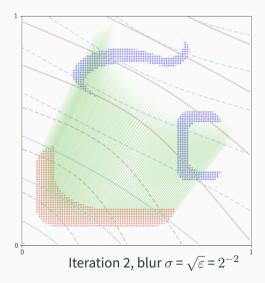
Visualizing F, G and the Brenier map $\nabla F(x_i) = -\frac{1}{\alpha_i}\partial_{x_i} \mathrm{OT}(\alpha,\beta)$

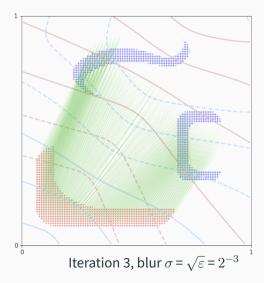


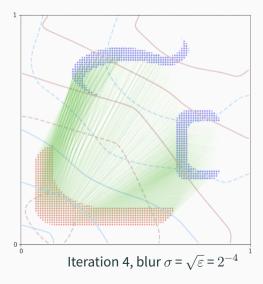
Visualizing F, G and the Brenier map $\nabla F(x_i) = -\frac{1}{\alpha_i}\partial_{x_i} \mathrm{OT}(\alpha,\beta)$

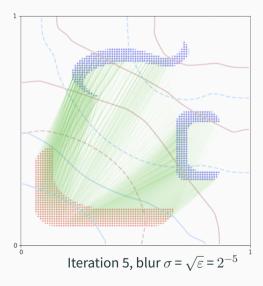


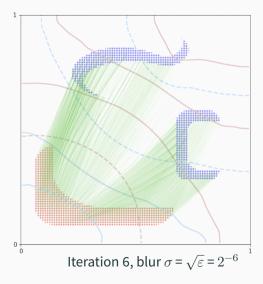
Visualizing F, G and the Brenier map $\nabla F(x_i) = -\frac{1}{\alpha_i}\partial_{x_i} \mathrm{OT}(\alpha,\beta)$

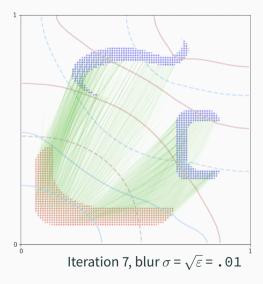


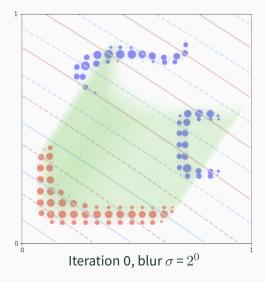


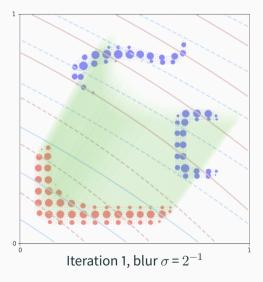


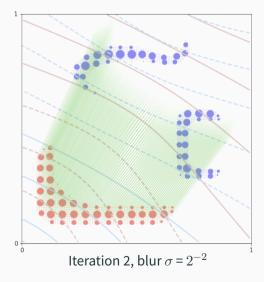


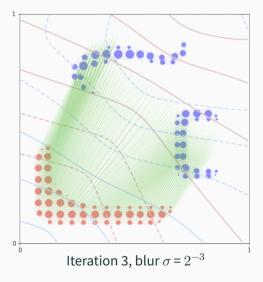


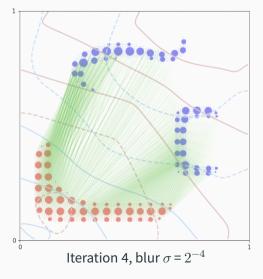


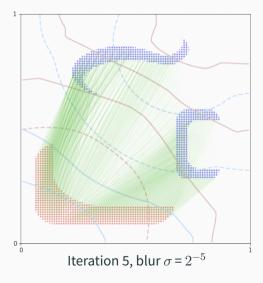


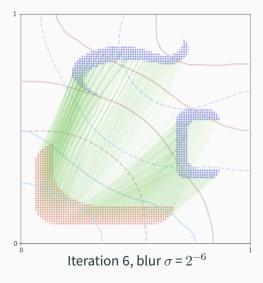


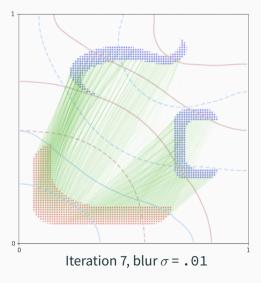










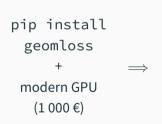


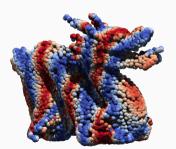
Scaling up optimal transport to anatomical data

Progresses of the last decade add up to a $\times 100 - \times 1000$ acceleration:

$$Sinkhorn~GPU \xrightarrow{\times 10} + KeOps \xrightarrow{\times 10} + Annealing \xrightarrow{\times 10} + Multi-scale$$

With a precision of 1%, on a modern gaming GPU:



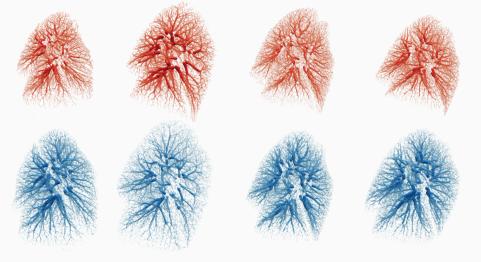


10k points in 30-50ms



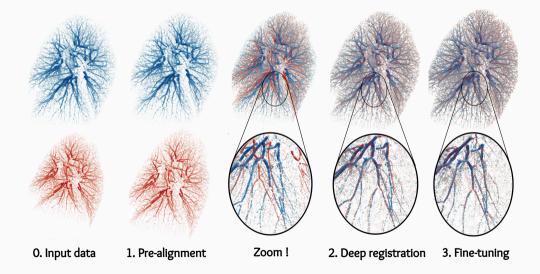
100k points in 100-200ms

A typical example in anatomy: lung registration "Exhale – Inhale"



 $\textbf{Complex} \ deformations, high \ \textbf{resolution} \ (50k-300k \ points), high \ \textbf{accuracy} \ (<1mm).$

Three-steps registration



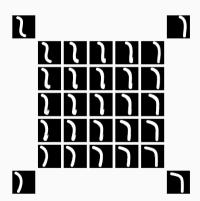
Wasserstein barycenters [AC11]

Barycenter
$$\mathbf{A}^* = \arg\min_{\mathbf{A}} \sum_{i=1}^{4} \lambda_i \operatorname{Loss}(\mathbf{A}, \mathbf{B}_i)$$
.



Euclidean barycenters.

$$\mathsf{Loss}(\mathsf{A},\mathsf{B}) = \|\mathsf{A} - \mathsf{B}\|_{L^2}^2$$



Wasserstein barycenters.

$$Loss(A, B) = OT(A, B)$$

Wasserstein barycenters

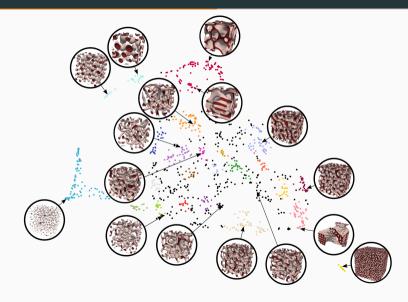
From a computational perspective:

- The problem is **convex** (easy) wrt. the weights.
- The support of the barycenter lies in the **convex hull** of the input distributions.

The curse of dimensionality hits hard:

- In high dimension, identifying the support can become **NP-hard**.
- In dimensions 2 and 3, we can just use a grid and recover super fast algorithms.
 Computing OT distances and barycenters between density maps is a solved problem.
- ⇒ We can now **easily** do manifold learning (= non-linear Model Order Reduction) in Wasserstein spaces of **2D and 3D** distributions.

An example: Anna Song's exploration of 3D shape textures [Son22]





Two very talented postdocs



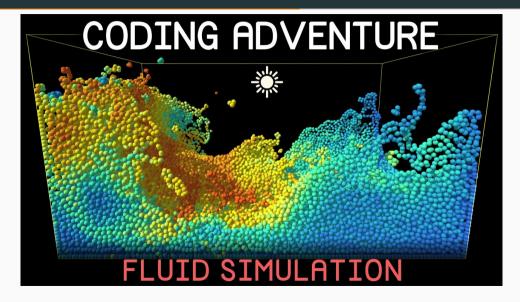
Maciej Buze Heriot-Watt University



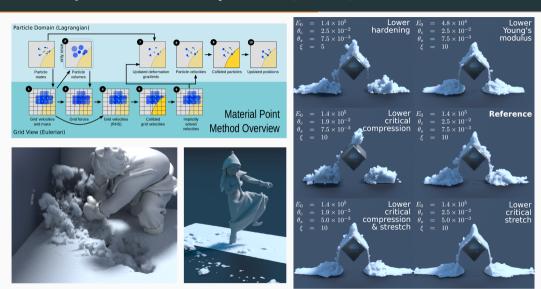
Antoine Diez Kyoto University

Original motivation: the N-body problem [Pri11]

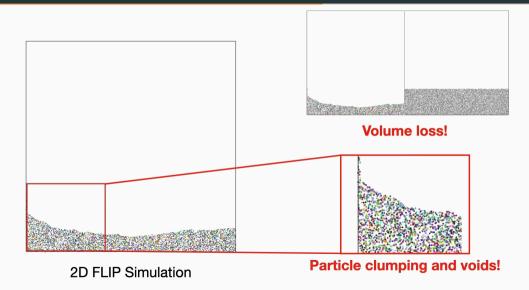




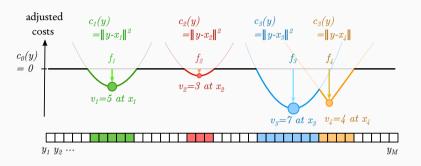
The material point method: Disney's Frozen [SSC+13]

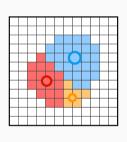


How can we enforce a volume preservation constraint? [QLDGJ22]



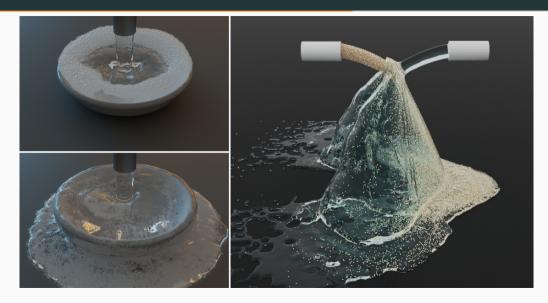
Use power diagrams i.e. semi-discrete optimal transport





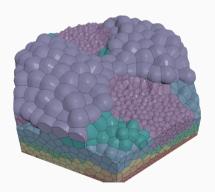
- The f_i 's maximize the dual objective $\sum_{i=1}^N v_i f_i + \int_{u \in \Omega} \min_{i=0}^N [\,c_i(y) f_i\,]\,\mathrm{d}y.$
- $\bullet \ \, \text{Optimality} \ \text{conditions} \ \iff \ \, \text{Vol}(\text{Cell}_i) = v_i.$
- To **compute the cells**, the objective and its gradient:
 - If $c_i(y) = \|y x_i\|^2$ for all cells, use a clever **grid-free** algorithm.
 - Otherwise, just use **KeOps**.

Power plastics [QLY⁺23]



Power plastics [QLY+23] – without the eye candy





Main numerical ingredients

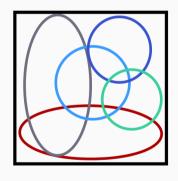
These simulations alternate between:

- 1. **Moving the particles** according to your favorite N-body model.
- 2. Computing Laguerre cells with the correct volumes:
 - (Multiscale) Sinkhorn for tolerance > 5%.
 - (Quasi-)Newton for tolerance < 1%.
- 3. **Correcting** the particle positions to enforce the volume-preservation constraint:
 - Jump to the centroid of the cell.
 - Or add a spring for smoother trajectories.

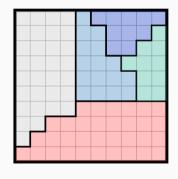
See e.g. Thomas Gallouët for a rigorous analysis with Mérigot, Lévy, etc.

But today: new applications with **custom cost functions** (thanks KeOps).

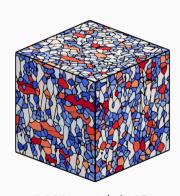
Anisotropic power diagrams let us model polycrystalline metals [BFR+24]



Ellipsoids.

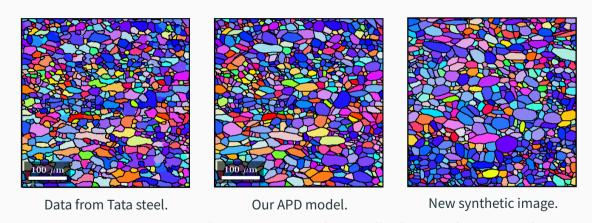


Pixel cells.



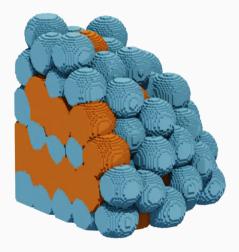
5,000 crystals in 3D.

Fit to real EBSD scan of low-carbon steel [BFR⁺24]

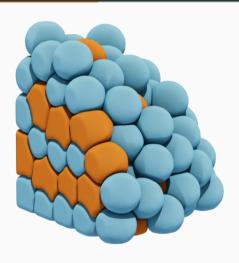


We can generate new, realistic 3D images with **prescribed properties** in seconds.

Change the cost function to simulate hard (blue) and soft (orange) cells [DF24]

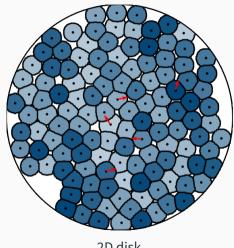


The **raw** 100x100x100 pixel grid...



with some Hollywood **makeup**.

Run-and-tumble motion [DF24]

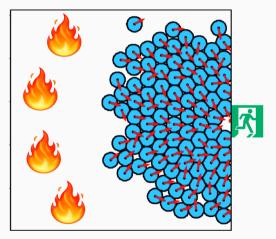


2D disk.

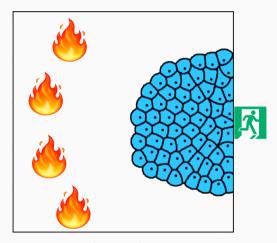


3D cube.

Fire alarm! [DF24]

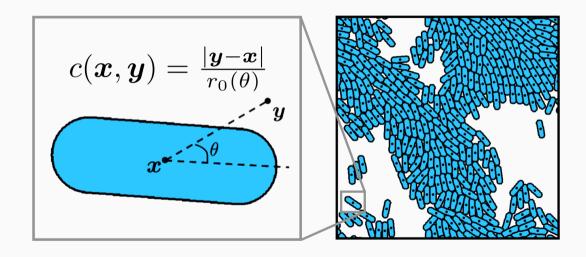


Hard particles **burn**.

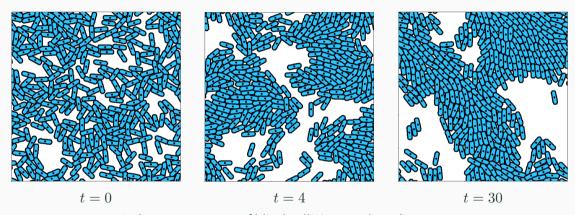


Soft particles **escape**.

Self-organizing swarms of blind, incompressible swimmers [DF24]

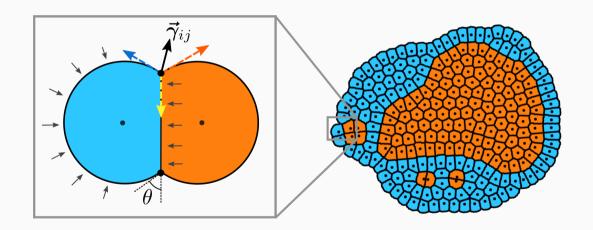


Self-organizing swarms of blind, incompressible swimmers [DF24]

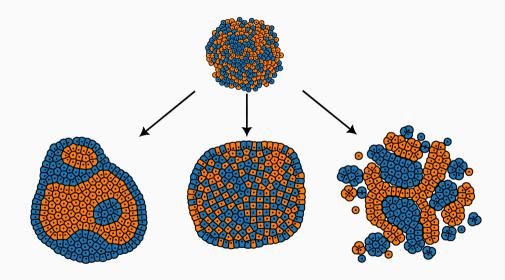


Order emerges out of blind collisions and re-alignments.

Surface tension [DF24]



Surface tension [DF24] – playing with the energy parameters





Conclusion

Genuine team work



Benjamin Charlier



Joan Glaunès



Thibault Séjourné



F.-X. Vialard



Gabriel Peyré



Alain Trouvé



Marc Niethammer



Shen Zhengyang



Olga Mula



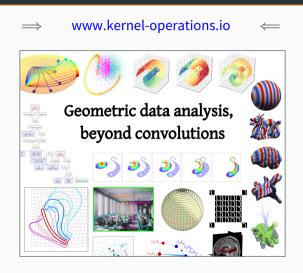
Hieu Do

Key points

- Optimal Transport = volume preservation = generalized sorting :
 - → Super-fast solvers on **simple domains**, especially 2D/3D spaces.
 - → **Fundamental tool** at the intersection of geometry and statistics.
- "Video-game physics" is great for modelling:
 - Expressive, real-time simulations that you can implement without being a Finite Elements guru: XPBD, DiffPD, Taichi...
- GPUs are more versatile than you think.
 - Ongoing work to provide **fast GPU backends** to researchers, going beyond what Google and Facebook are ready to pay for.

2026 target for scientific Python: **interactive**, **web-based** simulations à la ShaderToy.

Documentation and tutorials are available online



www.jeanfeydy.com/geometric_data_analysis.pdf

References

References i



M. Agueh and G. Carlier.

Barycenters in the Wasserstein space.

SIAM Journal on Mathematical Analysis, 43(2):904–924, 2011.



Dimitri P Bertsekas.

A distributed algorithm for the assignment problem.

Lab. for Information and Decision Systems Working Paper, M.I.T., Cambridge, MA, 1979.

References ii



Maciej Buze, Jean Feydy, Steven Roper, Karo Sedighiani, and David P Bourne.

Anisotropic power diagrams for polycrystal modelling: efficient generation of curved grains via optimal transport.

arXiv submission 5452163, 2024.



Haili Chui and Anand Rangarajan.

A new algorithm for non-rigid point matching.

In Computer Vision and Pattern Recognition, 2000. Proceedings. IEEE Conference on, volume 2, pages 44–51. IEEE, 2000.

References iii



Marco Cuturi.

Sinkhorn distances: Lightspeed computation of optimal transport.

In Advances in Neural Information Processing Systems, pages 2292–2300, 2013.



Antoine Diez and Jean Feydy.

An optimal transport model for dynamical shapes, collective motion and cellular aggregates, 2024.

References iv



Steven Gold, Anand Rangarajan, Chien-Ping Lu, Suguna Pappu, and Eric Mjolsness.

New algorithms for 2d and 3d point matching: Pose estimation and correspondence.

Pattern recognition, 31(8):1019–1031, 1998.



Leonid V Kantorovich.

On the translocation of masses.

In Dokl. Akad. Nauk. USSR (NS), volume 37, pages 199–201, 1942.

References v



Harold W Kuhn.

The Hungarian method for the assignment problem.

Naval research logistics quarterly, 2(1-2):83–97, 1955.



Jeffrey J Kosowsky and Alan L Yuille.

The invisible hand algorithm: Solving the assignment problem with statistical physics.

Neural networks, 7(3):477-490, 1994.

References vi



Sebastian Lague.

Coding adventure: Simulating fluids.

https://www.youtube.com/watch?v=rSKMYc1CQHE&t=1s, 2023.



Bruno Lévy.

A numerical algorithm for l2 semi-discrete optimal transport in 3d.

ESAIM: Mathematical Modelling and Numerical Analysis, 49(6):1693–1715, 2015.

References vii



Quentin Mérigot.

A multiscale approach to optimal transport.

In Computer Graphics Forum, volume 30, pages 1583–1592. Wiley Online Library, 2011.



Gabriel Peyré and Marco Cuturi.

Computational optimal transport.

arXiv preprint arXiv:1803.00567, 2018.

References viii



Anthony Prieur.

Simulation de la formation des structures de l'univers.

https://github.com/devpack/nbody-cosmos, 2011.



Ziyin Qu, Minchen Li, Fernando De Goes, and Chenfanfu Jiang.

The power particle-in-cell method.

ACM Transactions on Graphics, 41(4), 2022.

References ix



Ziyin Qu, Minchen Li, Yin Yang, Chenfanfu Jiang, and Fernando De Goes.

Power plastics: A hybrid Lagrangian/Eulerian solver for mesoscale inelastic flows.

ACM Transactions on Graphics (TOG), 42(6):1–11, 2023.



Bernhard Schmitzer.

Stabilized sparse scaling algorithms for entropy regularized transport problems.

SIAM Journal on Scientific Computing, 41(3):A1443–A1481, 2019.

References x



Anna Song.

Generation of tubular and membranous shape textures with curvature functionals.

Journal of Mathematical Imaging and Vision, 64(1):17–40, 2022.



Alexey Stomakhin, Craig Schroeder, Lawrence Chai, Joseph Teran, and Andrew Selle.

A material point method for snow simulation.

ACM Transactions on Graphics (TOG), 32(4):1–10, 2013.